Cryptography I, winter 2005/06 Joachim von zur Gathen, Michael Nüsken

10. Exercise sheet Hand in before Monday, 2006/01/30, 14⁰⁰ in b-it 1.22.

Exercise 10.1 (ElGamal signatures). (4 points)

We choose a prime number p=12347 and the group $G=\mathbb{Z}_p^{\times}$. We use a=9876 as the secret part of the key $K=(p,g,\beta,a)$. The message x to be signed consists of the last four digits of your student registration number. Use k=399 as your random number from $\mathbb{Z}_{p-1}^{\times}$.

- (i) Show: g = 2 generates G. Compute $\beta = g^a \in G$.
- (ii) Compute the signature $\operatorname{sig}_K(x,k) = (x,\gamma,\delta)$. Here $\gamma = g^k \bmod p$ and $\delta = (x-a\gamma)k^{-1} \bmod (p-1)$. Verify your signature.

Exercise 10.2 (A mysterious equation). (3+2 points)

Let $p \in \mathbb{N}$ be a prime number. The central operation in verifying an ElGamal signature is checking the equation $g^x = \beta^\gamma \cdot \gamma^\delta$ in \mathbb{Z}_p , where $g, \beta, \gamma \in \mathbb{Z}_p^\times$ and $\delta \in \mathbb{Z}_p$ and $x \in \{0, 1, \dots, p-1\}$ is the message or its hash, respectively. For now we consider the somewhat simpler congruence

$$(*) g^x = \beta^{\gamma} \cdot \gamma \quad \text{in } \mathbb{Z}_p$$

with $g, \beta \in \mathbb{Z}_p^{\times}$, $\gamma \in \mathbb{Z}_{p(p-1)}$ and $x \in \mathbb{N}$, $0 \le x \le p-1$.

- (i) Show: $\gamma = g^x(1-p)$ rem (p^2-p) is a solution to the equation (*).
- (ii) It holds that $\mathbb{Z}_{p(p-1)} \cong \mathbb{Z}_p \times \mathbb{Z}_{p-1}$. We identify $\mu \in \mathbb{Z}$, $0 \leq \mu < p$ with $(\mu,0) \in \mathbb{Z}_p \times \mathbb{Z}_{p-1}$. Let μ be a solution to congruence (*). For which $\ell \in \mathbb{Z}_{p-1}$ is there a $\lambda \in \mathbb{Z}_p$ so that also $(\mu \cdot \lambda, \ell) \in \mathbb{Z}_p \times \mathbb{Z}_{p-1}$ is a solution to (*)? Compute the dependency of λ on ℓ for that case.
- (iii) Is $\operatorname{sig}_K(x)=(x,g^x(1-p),1)$ a legal ElGamal signature? What is the consequence of this discovery for the practical use of the ElGamal signature scheme?
- (iv*) How many solutions γ are there for the congruence (*) and fixed g, β , x, +2 p?

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Exercise 10.3 (DSA Practice).

(10 points)

In this exercise you will make practical computations with the DSA algorithm, using *real life* key sizes.

- (i) Generate a random prime number q with exactly 160 bits.
- (ii) Generate a prime p with exactly 1024 bits, such that q divides p-1.
- (iii) Find a $g \in \mathbb{Z}_p^{\times}$ which has the exact order q. Let $G = \langle g \rangle \subset \mathbb{Z}_p^{\times}$ be the cyclic group with q elements generated by g.
- (iv) Let a < p be a random number and $y = g^a \in G$. We shall consider a to be Alice's secret key and y her public key.
- (v) Let $m \in \mathbb{Z}_q$ be the integer value of the ASCII text: DSA_for_real (note the two blanks in the text!). Using a random number $k \in \mathbb{Z}_q$ produce a DSA signature S(m) = (m, x, b) on the message m on behalf of Alice.
- (vi) Let Bob know the public key (p,g,y). Verify the signature S(m) on behalf of Bob.
- (vii) Let m' be the integer value of the ASCII text:

Can you produce a DSA signature of this text using the same setting as above? If no, what additional steps are required?

- (viii) The DSA system can be attacked in two different ways:
 - (a) By solving the index problem in the group G iwth q elements, with the baby-step giant step algorithm (or the Pollard- ϱ method) in this group.
 - (b) By solving the general discrete logarithm problem in \mathbb{Z}_p^{\times} , using the up to date Number Field Sieve. The complexity of this method is given by the function:

$$L(p) = \exp\left(1.992 \cdot \left(\log p \cdot (\log\log p)^2\right)^{1/3}\right).$$

The function log is the natural logarithm.

Compare the two estimated times, when $p \sim 2^{1024}$ and $q \sim 2^{160}$.