Security on the Internet, winter 2008

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5. Exercise sheet Hand in solutions until Monday, 08 December 2008, 11⁵⁹am (deadline!).

Note that on Wednesday, 3 December 2008 there is the Dies Academicus in Bonn and we will have no lecture/tutorial. We will thus on Tuesday, 2 December 2008 have a special session in the tutorial that repeats some of the major concepts regarding groups, rings, fields and other mathematical basics (or anything else you ask for).

As usual: Any claim needs a proof or an argument.

Exercise 5.1 (Exponentiation & discrete logarithms). (15+3 points)

Suppose *G* is a group and *g* is an element of order ℓ . In the course we have defined exponentiation as a map from the integers \mathbb{Z} to some group *G*.

(i) Show that it makes sense to view it as a map

$$\exp_g\colon \begin{array}{ccc} \mathbb{Z}_\ell & \longrightarrow & \langle g\rangle \subseteq G, \\ x & \longmapsto & g^x \end{array}.$$

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- (ii) Let $G = \mathbb{Z}_{10001}^{\times}$, g = 42. Write a procedure to compute \exp_g efficiently. 3 [Group operations are allowed as primitives. Other predefined procedures may not be used.]
- (iii) Same for $G = \mathbb{Z}_{241576501}^{\times}$, g = 23.
- (iv) Now let p = 241576501, and $g = 23^{1500} = -46436978 \in \mathbb{Z}_p^{\times}$.
 - (a) Compute g^{11^4} and g^{11^5} .
 - (b) Prove that the order of g is 11^5 .
 - (c) Prepare a table with all powers of $h := g^{11^4} = 23^{(p-1)/11}$ in \mathbb{Z}_p^{\times} .
 - (d) Compute the discrete logarithm x of $42^{1500} = 105868544 \in \mathbb{Z}_p^{\times}$ with respect to g. [Note that $(p-1) = 1500 \cdot 11^5$ and consider $42^{1500 \cdot 11^4} = g^{x \cdot 11^4} \cdots$]
 - (e) What does the result tell us about the discrete logarithm of $42 \in \mathbb{Z}_p^{\times}$ with ± 3 respect to the base $23 \in \mathbb{Z}_p^{\times}$?

Exercise 5.2 (High powers).	(3 points)
Compute $3^{98765432101}$ in \mathbb{Z}_{101} .	3

Exercise 5.3 (Pollard's *p* method).

(9 points)

In class we discussed Pollard's ρ method for computing the discrete logarithm in a group \mathbb{Z}_p^{\times} of size *m*. In particular we defined the algorithm in the following way: Assuming that we work on tuples $(\gamma, \delta, ag^{\gamma}, g^{\delta})$ we looked at some in a sense randomly behaving function *f* that mapped such tuples to other ones. This however is not efficient enough. [Why?] Thus we consider instead tuples $(\gamma, \delta, a^{\gamma}g^{\delta})$ and the function *f* defined as follows:

$$\begin{array}{cccc} \mathbb{Z}_{p-1} \times \mathbb{Z}_{p-1} \times \mathbb{Z}_{p}^{\times} & \longrightarrow & \mathbb{Z}_{p-1} \times \mathbb{Z}_{p-1} \times \mathbb{Z}_{p}^{\times}, \\ f: & & \begin{pmatrix} (\gamma, \delta, x) & \longmapsto & \begin{cases} (2\gamma, 2\delta, x^2) & x_1 = x_0 \\ (\gamma, \delta + 1, gx) & x_1 x_0 = 01 \\ (\gamma + 1, \delta, ax) & x_1 x_0 = 10 \end{cases}$$

- (i) We start at $(\gamma_0, \delta_0, a^{\gamma_0} g^{\delta_0})$ with $\gamma_0, \delta_0 \xleftarrow{\bullet} \mathbb{Z}_{p-1}$, and determine $(\gamma_i, \delta_i, x_i) = f^i(\gamma_0, \delta_0, a^{\gamma_0} g^{\delta_0})$. Show that $x_i = a^{\gamma_i} g^{\delta_i}$.
- (ii) Show that with a collision in the third coordinate one can easily compute the discrete logarithm of a to the base g.
- (iii) Show that this can be done with heuristically expected $O(\sqrt{m})$ group operations. You may assume that f indeed behaves randomly. Hint: Birthday-paradox.
- (iv) Implement Pollard's ρ algorithm and compute the discrete logarithm of your student registration number in the group \mathbb{Z}_p^{\times} with $p = 10^8 + 37$ and base g = 2. Count the number of group operations needed. Additionally hand in the source code.

Exercise 5.4 (Order).

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(8+6 points)

Let *G* be a (multiplicative) commutative group, *a* an element of order 24 and *b* an element of order 33. What is the order of b^2 , b^3 , ...? What are possible orders of *ab*?

Let us look at an example first: Take $G = \mathbb{Z}_{1321}^{\times}$, a = 17 and b = 53. We have $a^{24} = 1$ and $b^{33} = 1$ in G and for all respective smaller positive exponents the result is not 1.

(i) Compute the order of b^2 , b^3 , b^9 , b^{10} , b^{11} .

Now we want to investigate the general case:

(ii) Show: The order of the power g^k of a group element $g \in G$ is the order of g divided by the greatest common divisor of k and that order, in formulae: $\operatorname{ord}(g^k) = \operatorname{ord}(g)/\operatorname{gcd}(k, \operatorname{ord}(g))$. [*Hint*: Look at the special cases $\operatorname{gcd}(k, \operatorname{ord}(g)) = 1$ and $k | \operatorname{ord}(g)$ and derive the general solution from there.]

Consider again the example:

- (iii) Compute the order of ab, ab^2 , ab^3 .
- ... and back to the general case:
- (iv) Show: The order of the product xy of two group elements $g, h \in G$ in a commutative group G divides the least common multiple of the orders of g and h, in formulae: $\operatorname{ord}(gh) | \operatorname{lcm}(\operatorname{ord}(g), \operatorname{ord}(h))$.
- (v*) Show: If the orders of two group elements $g, h \in G$ are coprime, then the order of xy is actually equal to the the least common multiple of those orders, in formulae: $gcd(ord(g), ord(h)) = 1 \Rightarrow ord(gh) = lcm(ord(g), ord(h))$.
- (vi*) Show that the following is true in general: If $\operatorname{ord}(g) = ms$, $\operatorname{ord}(h) = mt + 3$ where *s* and *t* are coprime, then $st | \operatorname{ord}(gh)$.

[Actually, $st = \operatorname{lcm}(\operatorname{ord}(g), \operatorname{ord}(h)) / \operatorname{gcd}(\operatorname{ord}(g), \operatorname{ord}(h))$.]