

# Advanced Cryptography: Quantum Cryptography

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## 7. Exercise sheet

Hand in solutions until Monday, 13 December 2010, 23:59h.

**Exercise 7.1** (Phase estimation revisited).

(5+3 points)

In the lecture we discussed a quantum algorithm for phase estimation. There we wanted to approximate a phase  $\varphi$  using a  $t$ -bit binary number  $b/2^t$  (with  $b \in [0, 2^t - 1]$ ), where the error  $\delta = \varphi - b/2^t$  satisfied  $0 \leq \delta \leq 2^{-t}$ . We computed the amplitudes of  $|(b + \ell) \bmod 2^t\rangle$  as

$$\alpha_\ell = \frac{1}{2^t} \frac{1 - \exp(2\pi i(2^t \delta - \ell))}{1 - \exp(2\pi i(\delta - \ell/2^t))}$$

in case  $\delta \neq 0$  or  $\ell \neq 0$ . For  $\delta = 0$  we have  $\alpha_0 = 1$

(i) Set  $\varphi = 1/7$ . Plot the probabilities for the possible measurements for  $t \in \{3, 5, 8, 12\}$  at the points  $\ell/2^t$ , where  $\ell$  runs through  $[-2^{t-1}, 2^{t-1}] \cap \mathbb{Z}$ . 5

(ii) Interpret your results. +3

**Exercise 7.2** ( $\varphi$ -asco).

(3 points)

Let  $p, q \in \mathbb{N}$  be two different prime numbers. Let  $N = p \cdot q$ . Then  $\varphi(N) = (p - 1) \cdot (q - 1)$ , where  $\varphi$  is Euler's totient function. Compute — without factoring  $N$  —  $p$  and  $q$ , if  $N = 168149075693$  and  $\varphi(N) = 168148245408$  are given. *Hint*: Look at the quadratic equation  $x^2 - (p + q)x + pq = 0$ . 3

**Exercise 7.3** (Order mod  $N$ ).

(4 points)

(i) Compute the order of  $x = 5$  in  $\mathbb{Z}_{21}^\times$ . 1

(ii) Compute  $3^{1000003}$  in  $\mathbb{Z}_{101}^\times$  by hand. *Note*: Only a small calculation is needed! 3

**Exercise 7.4** (Another unitary operation).

(13 points)

Let  $N$  be a natural number,  $x \in \mathbb{Z}_N^\times$  and  $y \in \{0, 1\}^L$ . In the lecture we encountered the function

$$U|y\rangle = \begin{cases} |xy \bmod N\rangle & , \text{ if } 0 \leq y < N, \\ |y\rangle & , \text{ if } N \leq y < 2^L. \end{cases}$$

(i) Show that  $U$  is unitary. 3

(ii) Sketch the circuit for order finding. 2

In the lecture we computed the eigenstates  $|u_s\rangle$  of  $U$  as

$$|u_s\rangle = \frac{1}{\sqrt{r}} \sum_{0 \leq k < r} \exp\left(\frac{-2\pi i s k}{r}\right) |x^k \bmod N\rangle,$$

where  $r$  is the order of  $x$  in  $\mathbb{Z}_N$ .

3 (iii) Prove that  $\frac{1}{\sqrt{r}} \sum_{0 \leq s < r} |u_s\rangle = |1\rangle$ .

2 (iv) Show that in the algorithm the state before the inverse Fourier transform is

$$|\psi\rangle = \sum_{0 \leq j < 2^t} |j\rangle U^j |1\rangle = \sum_{0 \leq j < 2^t} |j\rangle |x^j \bmod N\rangle.$$

3 (v) Starting from this  $|\psi\rangle$  compute (again) the inverse Fourier transform of  $|\psi\rangle$  with respect to  $|j\rangle$ .

