# Advanced Cryptography: Algorithmic Cryptanalysis <br> Daniel Loebenberger, Konstantin Ziegler 

## 2. Exercise sheet

Hand in solutions until Saturday, 30 April 2011, 23:59h.

## Exercise 2.1 (Using an SPN for decryption).

Let $y$ be the encryption of a message $x$ with key $K$ by an SPN with S-box $\pi_{S}$ and bit-permutation $\pi_{P}$. In other words,

$$
y=\operatorname{SPN}\left(x, \pi_{S}, \pi_{P},\left(K^{1}, \ldots, K^{N+1}\right)\right),
$$

where ( $K^{1}, \ldots, K^{N+1}$ ) is the key schedule. Find an S-box $\pi_{S}^{*}$, a bit-permutation $\pi_{P}^{*}$ and a key schedule $\left(L^{1}, \ldots, L^{N+1}\right)$, such that

$$
x=\operatorname{SPN}\left(y, \pi_{S}^{*}, \pi_{P}^{*},\left(L^{N+1}, \ldots, L^{1}\right)\right) .
$$

## Exercise 2.2.

Suppose that the S-box of the example in the lecture is replaced by the S-box 7 defined by the following substitution:

| $z$ | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | A | B | C | D | E | F |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\pi_{S}(z)$ | E | 2 | 1 | 3 | D | 9 | 0 | 6 | F | 4 | 5 | A | 8 | C | 7 | B |

(i) Compute the table of values $N_{D}$ for this S-box.
(ii) Find a differential trail using four active S-boxes, namely $S_{1}^{1}, S_{4}^{1}, S_{4}^{2}$, and $S_{4}^{3}$, that has propagation ratio $27 / 2048$.
(iii) How many encrypted messages will you have to request for a differential attack with this trail in order to achieve similar confidence as with the differential trail described in the lecture?

## Exercise 2.3.

Suppose that $X_{1}, X_{2}$, and $X_{3}$ are independent discrete random variables defined on the set $\{0,1\}$. Let $\epsilon_{1}$ denote the bias of $X_{i}$, for $i=1,2,3$. Under which conditions on $\epsilon_{i}$ are $X_{1} \oplus X_{2}$ and $X_{2} \oplus X_{3}$ independent? (Recall, that in the lecture, we saw that this is in general not the case.)

## Exercise 2.4.

Daniel shows you his self-made random-number-generator which produces 16-bit numbers. But the distribution is not uniform! Daniel's favorite number is chosen with probability $27 / 1024$ - and you know that probability, but not the value of the number.
(i) How many calls to the random-number-generator do you expect to make, such that the favorite number occurs at least 9 times?
(ii) Assume that the probability distribution of the non-favorite numbers is uniform. What is the probability that any other number occurs at least 5 times, given the numbers of calls you derived in (i)?

