# Advanced Cryptography: Algorithmic Cryptanalysis <br> Daniel Loebenberger, Konstantin Ziegler 

## 9. Exercise sheet <br> Hand in solutions until Saturday, 18 June 2011, 23:59h.

To estimate the average effort you put into solving the following exercises, please add after each exercise the amount of time you spent for it.

Exercise 9.1 (Breaking truncated linear congruential generators).
(19+5 points)
We consider the truncated homogenous linear congruential generators with $x_{i}=s x_{i-1}$ rem $m$. We are given that $m=1009, k=\left\lceil\log _{2}(m) / 2\right\rceil=5$ and $a=25$. The sequence $y$ is defined as $y_{i}:=x_{i} \operatorname{div} 2^{k}$ which you intercepted as

$$
0,10,21,25,30,8,13,13,24,14,7,6,15,28,10,3,17,25,0,15,12, \ldots
$$

Our task is to break this generator completely. To do so, we will recover the sequence $z_{i}$ with $x_{i}=y_{i} 2^{k}+z_{i}$.
(i) Write down the matrix (over $\mathbb{Z}$ !)

$$
A=\left[\begin{array}{cccccc}
m & 0 & 0 & 0 & 0 & 0 \\
-s & 1 & 0 & 0 & 0 & 0 \\
-s^{2} & 0 & 1 & 0 & 0 & 0 \\
-s^{3} & 0 & 0 & 1 & 0 & 0 \\
-s^{4} & 0 & 0 & 0 & 1 & 0 \\
-s^{5} & 0 & 0 & 0 & 0 & 1
\end{array}\right]
$$

(ii) Compute the sequence $c_{i}:=\left(s^{i-1} y_{1}-y_{i}\right) 2^{k}$ over $\mathbb{Z}$ for $i=1, \ldots, 6$.
(iii) Using lattice basis reduction compute a reduced basis $B$ and a unimodular transformation $U$ such that $B=U A$.
(iv) Compute $U c$ and take the balanced system of representatives modulo $m$ of your result, i.e. divide each entry $c_{i}^{\prime}=(U c)_{i}$ with rest by $m$ and if it is larger than $m / 2$ write it as $c_{i}^{\prime}-m$.
(v) Now solve $B z=U c$ using Gaussian elimination, obtaining the $z_{i}$.
(vi) Finish by writing down the sequence $x_{i}$.
(vii) Compute the next 10 values of the above sequence of $y$ 's.
(viii) Argue that you have broken the generator.
(ix) Explain in detail why we had to use basis reduction at all.
(x) Play a bit around with your algorithms. Try different values of $m, s$ and $k$ and report on the successes and failures of your algorithm.

Exercise 9.2 ( $\varphi$-asco).

Exercise 9.3 (Factoring $x$ modulo $N$ ).
Let $p \neq q$ be prime numbers, $N=p \cdot q, f=x \in \mathbb{Z}_{N}[x]$.
(i) Show that $p^{2}+q^{2}$ is a unit in $\mathbb{Z}_{N}^{\times}$, i.e. $\operatorname{gcd}\left(p^{2}+q^{2}, p q\right)=1$.
(ii) Let $u \in \mathbb{Z}_{N}$ be the inverse of $p^{2}+q^{2}$. Show that $f=u(p x+q)(q x+p)$.
(iii) Prove that the two linear factors in (ii) cannot be written as a product of two polynomials of degree at least 1 and that they cannot be inverted in $\mathbb{Z}_{N}[x]$. Hint: Consider the situation in $\mathbb{Z}_{p}$ and and $\mathbb{Z}_{q}$ separately.
(iv) Conclude that if you can factor polynomials over $\mathbb{Z}_{N}$ then you can factor the integer $N$.

## Exercise 9.4 (A lattice basis).

Let $e, d, p, q, N$ be positive natural numbers defines as in the specification of the RSA cryptosystem. In the lecture we discussed a lattice whose elements are the exactly the pairs $\left(x_{1}, x_{2}\right) \in \mathbb{Z}^{2}$ of solutions to the equation $e x_{1}+x_{2}=0$ in $\mathbb{Z}_{N}$. In this exercise we will explore a more general problem: Let $a_{1}, \ldots, a_{n} \in \mathbb{Z}$ and $N \in \mathbb{N}$ with $\operatorname{gcd}\left(a_{n}, N\right)=1$ and consider the lattice

$$
L=\left\{\left(x_{1}, \ldots, x_{n}\right) \in \mathbb{Z}^{n} \mid a_{1} x_{1}+\cdots+a_{n} x_{n}=0 \text { in } \mathbb{Z}_{N}\right\}
$$

(ii) Rewrite the resulting equation in $\mathbb{Z}_{N}$ as an equation over the integers.
(iii) Find a basis for $L$.
(iv) Describe how you can use the above lattice to find solutions for the following problem: Input is a positive integer $N$ and integers $a_{1}, \ldots, a_{n}$. Find $x_{1}, \ldots, x_{n} \in\{0,1\}$ such that $\sum_{1 \leq i \leq n} a_{i} x_{i}=0$ in $\mathbb{Z}_{N}$.

